

Enriched Categories and Applications

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1 Motivation

2 Definition

- Monoidal Category
- Closed
- Enriched Category

3 Example

4 Application

Motivation

For objects $A, B \in \mathcal{C}$, we denote $\mathcal{C}(A, B)$ as the collection of morphisms from A to B . Two interesting question:

1. Given a category \mathcal{C} , what is the structure of such collections, if any?

Example

- For a locally small category \mathcal{C} , the collection $\mathcal{C}(A, B)$ is a set for all objects $A, B \in \mathcal{C}$. Therefore, we can denote it $S = \mathcal{C}(A, B) \in \mathbf{Set}$.
- In a pre-additive category \mathcal{C} , the collection $\mathcal{C}(A, B)$ always has a structure of abelian group such that the composition is bilinear for all objects $A, B \in \mathcal{C}$. Therefore, we can denote it $G = \mathcal{C}(A, B) \in \mathbf{Ab}$.

Enrichment: $\mathcal{C}(A, B)$ “looks like” objects from some category.

2. When is a category enriched over itself?

Example

Ab, **Vect** $_{\mathbb{F}}$ and **R-Mod** are categories enriched over themselves.

Definition (Monoidal Category)

A **monoidal category** $\mathcal{V} = (\mathcal{V}_0, \otimes, I, a, l, r)$ consists of a category \mathcal{V}_0 , a functor $\otimes : \mathcal{V}_0 \times \mathcal{V}_0 \rightarrow \mathcal{V}_0$, an object I of \mathcal{V}_0 and natural isomorphisms $a_{XYZ} : (X \otimes Y) \otimes Z \rightarrow X \otimes (Y \otimes Z)$, $l_X : I \otimes X \rightarrow X$, $r_X : X \otimes I \rightarrow X$, that satisfies two coherence axioms expressed by the diagram below:

$$\begin{array}{ccc} ((W \otimes X) \otimes Y) \otimes Z & \xrightarrow{a} & (W \otimes X) \otimes (Y \otimes Z) \xrightarrow{a} W \otimes (X \otimes (Y \otimes Z)) \\ \downarrow a \otimes 1 & & \uparrow 1 \otimes a \\ (W \otimes (X \otimes Y)) \otimes Z & \xrightarrow{a} & W \otimes ((X \otimes Y) \otimes Z) \end{array}$$

Figure: Pentagon Rule

$$\begin{array}{ccc} (X \otimes I) \otimes Y & \xrightarrow{a} & X \otimes (I \otimes Y) \\ & \searrow r \otimes 1 & \swarrow 1 \otimes l \\ & X \otimes Y & \end{array}$$

Figure: Triangle Rule

Definition (Closed)

A monoidal category \mathcal{V} is **closed** if each functor $- \otimes Y : \mathcal{V}_0 \rightarrow \mathcal{V}_0$ has a right adjoint $[Y, -]$. In particular, we have the adjunction

$$\pi : \mathcal{V}_0(X \otimes Y, Z) \cong \mathcal{V}_0(X, [Y, Z])$$

or in terms of the unit-counit adjunction with unit $d : X \rightarrow [Y, X \otimes Y]$ and counit $e : [Y, Z] \otimes Y \rightarrow Z$. Notice that when $X = I$, by the isomorphism $l : I \otimes Y \cong Y$, we get a natural isomorphism

$$\mathcal{V}_0(Y, Z) \cong V[Y, Z]$$

where V is the functor $\mathcal{V}_0(I, -) : \mathcal{V}_0 \rightarrow \mathbf{Set}$. Therefore, the object $[Y, Z]$ acts as a lift of the hom set $\mathcal{V}_0(Y, Z)$, and is called the **internal hom**.

Remark

Just like in basic category theory, we focus our attention on locally small monoidal categories.

Definition (\mathcal{V} -enriched Category)

Let \mathcal{V} be a monoidal category, then a \mathcal{V} -**enriched category** \mathcal{A} consists of a set of objects $\text{ob}(\mathcal{A})$, a hom-object $\mathcal{A}(A, B) \in \mathcal{V}_0$ for each pair of objects of \mathcal{A} , a composition law $M = M_{ABC} : \mathcal{A}(B, C) \otimes \mathcal{A}(A, B) \rightarrow \mathcal{A}(A, C)$ for each triple of objects, and an identity element $j_A : I \rightarrow \mathcal{A}(A, A)$ for each object; this is subject to the associativity and unit axioms expressed by the commutativity of two diagrams:

$$\begin{array}{ccc}
 (\mathcal{A}(C, D) \otimes \mathcal{A}(B, C)) \otimes \mathcal{A}(A, B) & \xrightarrow{a} & \mathcal{A}(C, D) \otimes (\mathcal{A}(B, C) \otimes \mathcal{A}(A, B)) \\
 \downarrow M \otimes 1 & & \downarrow 1 \otimes M \\
 \mathcal{A}(B, D) \otimes \mathcal{A}(A, B) & & \mathcal{A}(C, D) \otimes \mathcal{A}(A, C) \\
 \searrow M & & \swarrow M \\
 & \mathcal{A}(A, D) &
 \end{array}$$

Figure: Associativity on Enriched Category

and

$$\begin{array}{ccccc}
 \mathcal{A}(B, B) \otimes \mathcal{A}(A, B) & \xrightarrow{M} & \mathcal{A}(A, B) & \xleftarrow{M} & \mathcal{A}(A, B) \otimes \mathcal{A}(A, A) \\
 \uparrow j \otimes 1 & \nearrow l & & \nwarrow r & \uparrow 1 \otimes j \\
 I \otimes \mathcal{A}(A, B) & & & & \mathcal{A}(A, B) \otimes I
 \end{array}$$

Figure: Unitality on Enriched Category

Examples

- Any pre-additive category can be enriched over **Ab**. In particular, abelian categories are enriched over **Ab**. Moreover, **Ab** is enriched over itself.
- Any (strict) 2-category can be enriched over **Cat**, the category of small categories.
- Any poset category can be enriched over the category of truth values, denoted $\mathcal{2}$.



Theorem: Monoidal closed categories are self-enriched

In particular, the structure of $\mathcal{V}(A, B)$ given a monoidal closed category \mathcal{V} and $A, B \in \mathcal{V}$ is $\mathcal{V}(A, B) = [A, B]$, which is the internal hom.

Remark

*The converse does not hold in general: observe that **Top** is a monoidal category that is not closed (the product functor $- \otimes Y$ does not commute with colimits, therefore does not have a right adjoint), but it can be self-enriched regardless: consider the structure **Top**(A, B) as equipped with the trivial topology where every subset is open.*

Example: $([0, 1], \leq)$

- Consider $([0, 1], \leq)$ as a monoidal (poset) category, with \leq as the usual comparison relation. Again, this is enriched over the category of truth values $\mathbb{2}$.
- The categorical product \times is given by the minimum function $\min(X, Y)$. The terminal object of the poset category is 1. By taking $\otimes = \times$ and $I = 1$, we obtain a (cartesian) monoidal category.
- Is this category closed? i.e. for $X, Y, Z \in [0, 1]$, is there a hom object structure $[Y, Z]$ such that there is a natural isomorphism $\mathbf{Hom}(X \otimes Y, Z) \cong \mathbf{Hom}(X, [Y, Z])$?

Yes, denote $[Y, Z] = \begin{cases} 1, & \text{if } Y \leq Z \\ Z, & \text{if } Y > Z \end{cases}$. This is the Gödel conditional, an operator of Gödel logics.

Example: still $([0, 1], \leq)$, but less canonical

- This time, denote \otimes as the usual multiplication operation $\times : [0, 1] \times [0, 1] \rightarrow [0, 1]$ on numbers. This is not a categorical product.
- Still take $I = 1$.
- The category is still closed. For $X, Y, Z \in [0, 1]$, there is a hom object structure $[Y, Z]$ such that $\mathbf{Hom}(X \otimes Y, Z) \cong \mathbf{Hom}(X, [Y, Z])$. This time, $[Y, Z] = \begin{cases} 1, & \text{if } Y \leq Z \\ \frac{Z}{Y}, & \text{if } Y > Z \end{cases}$. This is “the product structure”.

Application: Fuzzy Logic

Suppose a group of students are taking Putnam. Suppose A, B and C scores 10, 50 and 100 respectively. We may consider a fuzzy preorder \preccurlyeq as “at most as good as”.

What can we say about the relative performance between the three individuals using fuzzy logic? Intuition: $C \preccurlyeq A$ should be very close to 0, $B \preccurlyeq A$ should be less closer to 0, $C \preccurlyeq B$ should be much closer to 1.

Let $\mathcal{C} = (A, B, C)$ form a preorder category with respect to \preccurlyeq , then the category can be considered as “enriched” over $([0, 1], \leq)$ with

$$\mathcal{C}(X, Y) = \begin{cases} 1, & \text{if } X \leq Y \\ \frac{\text{score}(Y)}{\text{score}(X)}, & \text{if } X > Y \end{cases}. \text{ In particular, we have}$$

$$\mathcal{C}(C, A) = \frac{10}{100} = 0.1, \mathcal{C}(B, A) = \frac{10}{50} = 0.2 \text{ and } \mathcal{C}(C, B) = \frac{50}{100} = 0.5,$$

This can be quantified (“enriched”) over the category $([0, 1], \leq)$.

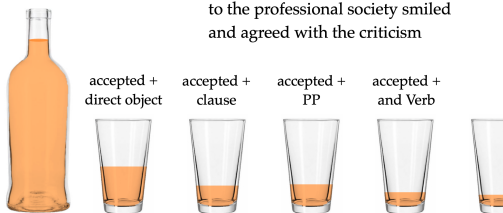
Application: Language Model

From the perspective of psycholinguistics,

- language prediction starts almost immediately after the speech starts!
i.e. **we start to predict what others are going to say as soon as they start talking.**
- this prediction is **ranked parallel** instead of all-or-none, i.e. our brain evaluates the most possible follow-up dynamically instead of just picking the suggestion with the highest probability.

The photographer accepted

the criticism
(that) her agent disagreed with her
to the professional society smiled
and agreed with the criticism



Levy (2008), *Cognition*

Application: Language Model

Bradley and Vlassopoulos (2021) described “a way of assigning to any word or expression s in a language a particular linear operator ρ_s .” From the perspective of the Yoneda lemma, the “meaning” of the word s is described by ρ_s , as the operator captures the environment and hence the meanings of word/expression.

Finally, they showed that the assignment $s \mapsto \rho_s$ is a functor between categories enriched over probabilities, i.e. $([0, 1], \leq)$. More precisely, this is an enriched functor from language (as a preorder category enriched over probabilities) to the category of those associated operators enriched over probabilities.



Max Kelly (2005)

Basic Concepts of Enriched Category Theory

Reprints in Theory and Applications of Categories, No. 10, 2005.



Tai-Danae Bradley and Yiannis Vlassopoulos (2021)

Language Modeling with Reduced Densities

Compositionality 3, 4 (2021).



Tai-Danae Bradley (2020)

Language Modeling with Reduced Densities (Blog Post)

<https://www.math3ma.com/blog/language-modeling-with-reduced-densities>